

The Nullstellensatz for algebraic structures II

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Nullstellensatz for finite fields

Let F_q be a finite field of order $q = p^n$.

Nullstellensatz for F_q

Let $\bar{x} = (x_1, \dots, x_n)$, and let $S : \{f_1(\bar{x}), \dots, f_m(\bar{x})\}$ be a system of equations over F_q . Then

$$\text{Rad}(S) = \text{id}\langle f_1, \dots, f_m, x_1^q - x_1, \dots, x_n^q - x_n \rangle$$

1 The Hilbert Nullstellensatz

2 The Nullstellensatz for fields

- The real numbers field
- The p -adic numbers field

3 Problems of the Nullstellensatz type

- The Resolvent Problem
- Statement of the Nullstellensatz problem
- The Nullstellensatz procedure for equationally Noetherian groups
- The irreducible Nullstellensatz problem

4 The Nullstellensatz for particular algebraic structures

- Torsion-free abelian groups
- Free metabelian groups

5 Questions and problems

The resolvent problem

Example: quadratic equations over \mathbb{R}

$s : ax^2 + bx + c = 0$, where a, b, c are parameters.

Let T be a transformation so that $T(s) = s'$,

$s' : y^2 = d$, where d is the discriminant.

Then T induces an isomorphism

$$T^* : \mathbb{R}[x]/\text{Rad}(s) \rightarrow \mathbb{R}[x]/\text{Rad}(s')$$

$R = \{\text{the set of resolvents for } s\} =$

$$\{r_0 : 0 = 1; r_1 : x = \sqrt{d}; r_2 : x = -\sqrt{d}\},$$

note that $c(r_i) < c(s)$.

Example: quadratic equations over \mathbb{R}

$d < 0$ Then s is inconsistent, $\text{Rad}(s) = R[x]$. The corresponding universal formula is

$$\Phi(a, b, c) : \forall x(ax^2 + bx + c = 0 \rightarrow 0 = 1).$$

This formula holds in \mathbb{R} iff $d = b^2 - 4ac < 0$.

$d \geq 0$

$$\Phi(a, b, c) : \forall x(ax^2 + bx + c = 0 \rightarrow x = \sqrt{d} \vee x = -\sqrt{d}).$$

We have

$$V(S) \subseteq V(r_1) \cup V(r_2).$$

Formal definition of the resolvent method

Let S be a finite system of equations with parameters over \mathcal{A} (over a field) in variables $X = \{x_1, \dots, x_n\}$.

Let $c(S)$ be complexity of systems of equations. We will say that for the system S there is a resolvent R if

- 1** $R = \{R_1, \dots, R_m\}$, $c(R_i) < c(S)$;

- 2** $V(R) \subseteq \bigcup_{i=1}^m R_i$; There exists finitely many universal formulas

$$\phi_i : \forall x(S(X) = 0 \rightarrow \bigvee_{j_1 \in J \subseteq \{1, \dots, m\}} R_{j_1}(X),$$

...

$$\phi_k : \forall x(S(X) = 0 \rightarrow \bigvee_{j_k \in J \subseteq \{1, \dots, m\}} R_{j_k}(X).$$

If such R exists, then we say that S has a solution of the resolvent problem.

Similar for systems of equations with parameters.

Definable predicate for inconsistent systems

Let $S(X, p_1, \dots, p_m)$ be a finite system of equations with parameters

$$\Phi(p_1, \dots, p_m) = \forall \bar{x} S(\bar{x}, p_1, \dots, p_m) = 0 \rightarrow 0 = 1)$$

Note that

$\Phi(p_1, \dots, p_m)$ is true $\Leftrightarrow S(\bar{x}, p_1, \dots, p_m) = 0$ is inconsistent.






Questions and open problems

- 1** The resolvent problem for systems of equations in one variable for finitely generated t.f. nilpotent groups.
- 2** Romanovskii and R.: solution of the resolvent problem for one-variable equations over some finitely generated metabelian groups.
- 3** The resolvent problem for free associative, Lie and anti-commutative algebras.

Let $\mathcal{A} = \langle A \mid L \rangle$ be an algebraic structure in a language L .

Problem of the Nullstellensatz type

For a system of equations S over \mathcal{A} find the radical $\text{Rad}(S)$.

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Algebraic geometry over groups I: Algebraic sets and ideal theory
Journal of Algebra, **219**, 16–79, 1999.
-  A. Myasnikov, V. Remeslennikov
Algebraic geometry over groups II: Logical foundations
Journal of Algebra, **234**, 225–276, 2000.
-  B. I. Plotkin
Some notions of algebraic geometry in universal algebra
Algebra and Analysis, **9** (4), 224–248, 1997.
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E. Daniyarova, A. Myasnikov, V. Remeslennikov

Unification theorems in algebraic geometry

Algebra and Discrete Mathematics, **1**, 80–112, 2008,

arXiv:0808.2522v1 [math.AG].



E. Daniyarova, A. Myasnikov, V. Remeslennikov

Algebraic geometry over algebraic structures II: Foundations

arXiv:1002.3562v2 [math.AG].

Unification Theorem A

Let \mathcal{A} be an equationally Noetherian algebra in a language L (with no predicates). Then for a finitely generated algebra \mathcal{C} of L the following conditions are equivalent:

- 1 $\text{Th}_{\forall}(\mathcal{A}) \subseteq \text{Th}_{\forall}(\mathcal{C})$, i.e., $\mathcal{C} \in \mathbf{Ucl}(\mathcal{A})$;
- 2 $\text{Th}_{\exists}(\mathcal{A}) \supseteq \text{Th}_{\exists}(\mathcal{C})$;
- 3 \mathcal{C} embeds into an ultrapower of \mathcal{A} ;
- 4 \mathcal{C} is discriminated by \mathcal{A} ;
- 5 \mathcal{C} is a limit algebra over \mathcal{A} ;
- 6 \mathcal{C} is an algebra defined by a complete atomic type in the theory $\text{Th}_{\forall}(\mathcal{A})$ in L ;
- 7 \mathcal{C} is the coordinate algebra of an **irreducible** algebraic set over \mathcal{A} defined by a system of equations in the language L .

Unification Theorem C

Let \mathcal{A} be an equationally Noetherian algebra in a language L (with no predicates). Then for a finitely generated algebra \mathcal{C} of L the following conditions are equivalent:

- 1 $\mathcal{C} \in \mathbf{Qvar}(\mathcal{A})$, i.e., $\text{Th}_{\text{qi}}(\mathcal{A}) \subseteq \text{Th}_{\text{qi}}(\mathcal{C})$;
- 2 $\mathcal{C} \in \mathbf{Pvar}(\mathcal{A})$;
- 3 \mathcal{C} embeds into a direct power of \mathcal{A} ;
- 4 \mathcal{C} is separated by \mathcal{A} ;
- 5 \mathcal{C} is a subdirect product of finitely many limit algebras over \mathcal{A} ;
- 6 \mathcal{C} is an algebra defined by a complete atomic type in the theory $\text{Th}_{\text{qi}}(\mathcal{A})$ in L ;
- 7 \mathcal{C} is the coordinate algebra of an algebraic set over \mathcal{A} defined by a system of equations in the language L .

The Nullstellensatz for algebraic structures

Terms and atomic formulas

Let $\mathcal{A} = \langle A \mid L \rangle$ be an algebraic structure in a functional language

$$L = \{\text{set of const symbols}\} \cup \{\text{set of functional symbols}\}.$$

Terms in the language L in variables $X = \{x_1, \dots, x_n\}$ are formal expressions defined recursively as follows:

- variables x_1, x_2, \dots, x_n and constants from L are terms;
- if t_1, \dots, t_n are terms and $F(x_1, \dots, x_n) \in L$ is a function then $F(t_1, \dots, t_n)$ is a term.

Atomic formulas are formulas of the form

$$(t = s),$$

where t, s are terms.

Denote by $At_L(X)$ the set of all atomic formulas in the language L with variables in X .

Elements of algebraic geometry

- Any subset $S \subseteq \text{At}_L(X)$ is called a **system of equations** in the language L .
- The set

$$V(S) = \{ (a_1, \dots, a_n) \in A^n \mid \\ t(a_1, \dots, a_n) = s(a_1, \dots, a_n) \quad \forall (t = s) \in S \},$$

is called **algebraic set** over \mathcal{A} .

- The **radical** $\text{Rad}(S)$ of S over \mathcal{A} is the set of all consequences of S over \mathcal{A} .
- An equation $(t' = s') \in \text{At}_L$ is called a **consequence** of S over \mathcal{A} if \mathcal{A} satisfies the (infinite) quasi-identity

$$\forall x_1, \dots, x_n \left(\bigwedge_{(t=s) \in S} t(\bar{x}) = s(\bar{x}) \longrightarrow t'(\bar{x}) = s'(\bar{x}) \right).$$

The Nullstellensatz for algebraic structure \mathcal{A} is ...

- ... a procedure for constructing the $\text{Rad}(S)$ of a system of equations S ;
- ... a theorem about the structure of the radical $\text{Rad}(S)$ of a system of equations S ;
- ... a procedure for constructing a system of equations S' with the Nullstellensatz property and so that $V(S') = V(S)$.

Definition

A system of equations S' is called a **system with Nullstellensatz property** if $\text{Rad}(S') = [S']$, where $[S']$ is the congruence closure of S' (e.g., normal subgroup, ideal).

Equationally Noetherian algebraic structures

Definition

An algebraic structure \mathcal{A} is called **equationally Noetherian** if for any positive integer n and any system of equations $S(x_1, \dots, x_n)$ there exists a finite subsystem $S_0 \subseteq S$ such that $V(S) = V(S_0)$.

Equationally Noetherian = every chain

$$\text{Rad}(S_1) \subseteq \text{Rad}(S_2) \subseteq \dots \subseteq \text{Rad}(S_m) \subseteq \dots$$

stabilizes.

Definition

An algebraic structure \mathcal{A} is called **Noetherian** (with respect to congruences) if ...

The Nullstellensatz for equationally Noetherian algebraic structures

Suppose that \mathcal{A} is an equationally Noetherian algebraic structure and S a system of equations.

Then $\text{Rad}(S)$ = the set of all equations $(t' = s')$ such that the quasi-identity

$$\forall x_1, \dots, x_n \left(\bigwedge_{(t=s) \in S} t(\bar{x}) = s(\bar{x}) \longrightarrow t'(\bar{x}) = s'(\bar{x}) \right)$$

holds in \mathcal{A} .

But, as usual, we don't know all quasi-identities that hold in \mathcal{A} .

Suppose, we know only axioms of $\mathbf{Qvar}(\mathcal{A})$. How can we realize the Nullstellensatz procedure in this case?

The Nullstellensatz procedure for equationally Noetherian groups

Let G be a group and Ω a “good” set of axioms for $\mathbf{Qvar}(G)$.

$$G[X] = G * F\langle x_1, \dots, x_n \rangle.$$

Equations = expressions of the type $s = 1$, where $s \in G[X]$.

To solve the Nullstellensatz problem it suffices to compute radicals in $G[X]$.

The Nullstellensatz procedure for equationally Noetherian groups

A procedure for constructing $\text{Rad}(S)$, S - finite:

$$S_0 = S \subseteq S_1 = \text{ncl}\langle S_0 \rangle \subset \dots \subset S_{i-1} \subset S_i \subset \dots$$

Here:

- $S_i = \text{ncl}\langle S_{i-1} \cup h_i \rangle$, $h_i \notin S_{i-1}$;
- $h_i = \mathbf{s}(w_1(x_1, \dots, x_n), \dots, w_r(x_1, \dots, x_n))$, where
- $w_1(x_1, \dots, x_n), \dots, w_r(x_1, \dots, x_n) \in G[X]$ and
- there exists a quasi-identity from Ω

$$\forall x_1, \dots, x_r \left(\bigwedge_{j=1}^m \mathbf{s}_j(\bar{x}) = 1 \longrightarrow \mathbf{s}(\bar{x}) = 1 \right);$$

- such that $\mathbf{s}_j(w_1(x_1, \dots, x_n), \dots, w_r(x_1, \dots, x_n)) \in S_{i-1}$ for all $j = 1, \dots, m$.

We have $\text{Rad}(S) = \bigcup_i S_i$. If G is **Noetherian** wrt congruences (normal subgroups), then $\text{Rad}(S) = S_i$ for some $i \in \mathbb{N}$.

The Nullstellensatz procedure for equationally Noetherian groups

A procedure for constructing the radical $\text{Rad}(S)$ for a finite system of equations S :

$$S_0 = S \subseteq S_1 = \text{ncl}\langle S_0 \rangle \subset \dots \subset S_{i-1} \subset S_i \subset \dots$$

Here:

- $S_i = \text{ncl}\langle S_{i-1} \cup \{\text{all } h_i\} \rangle$, $h_i \notin S_{i-1}$;
- $h_i = \mathbf{s}(w_1(x_1, \dots, x_n), \dots, w_r(x_1, \dots, x_n))$, where
- $w_1(x_1, \dots, x_n), \dots, w_r(x_1, \dots, x_n) \in G[X]$ and
- there exists a quasi-identity from Ω

$$\forall x_1, \dots, x_r \left(\bigwedge_{j=1}^m \mathbf{s}_j(\bar{x}) = 1 \longrightarrow \mathbf{s}(\bar{x}) = 1 \right);$$

- such that $\mathbf{s}_j(w_1(x_1, \dots, x_n), \dots, w_r(x_1, \dots, x_n)) \in S_{i-1}$ for all $j = 1, \dots, m$.

If G is **equationally Noetherian**, then $\text{Rad}(S) = S_i$ for some $i \in \mathbb{N}$.

Irreducible algebraic sets

Fact

Every algebraic set Y over an equationally Noetherian algebraic structure \mathcal{A} can be expressed as a finite union of irreducible algebraic sets (**irreducible components**):

$$Y = Y_1 \cup Y_2 \cup \dots \cup Y_m.$$

Furthermore, this decomposition is unique up to permutation and omission of superfluous irreducible components.

Irreducible = can not be expressed as a finite union of proper algebraic subsets.

Statement of the problem

The irreducible Nullstellensatz:

For a system of equations S

- find systems of equations S_1, \dots, S_m , such that

$$V(S) = V(S_1) \cup \dots \cup V(S_m)$$

and $V(S_i)$'s are irreducible,

- and compute the radicals $\text{Rad}(S_1), \dots, \text{Rad}(S_m)$.

Let G be an equationally Noetherian group and $S(X) \subseteq G * F(X)$ a finite system of equations over G . Suppose that we have a good system of axioms for $\mathbf{Ucl}(G)$ by formulas of the form:

$$\forall x_1, \dots, x_n \left(\bigwedge_{i=1}^m u_i(X) = 1 \rightarrow \bigvee_{j=1}^r v_j(X) = 1 \right)$$

Let $S_1 = \text{ncl}\langle S \rangle \triangleleft G[X]$. If $S_1 = \text{Rad}(S)$, then the process ends and we say that S has Nullstellensatz.

It is well-known that majority of quadratic equations have Nullstellensatz.

The Nullstellensatz

for torsion-free abelian groups



A. Myasnikov, V. Remeslennikov

Algebraic geometry over groups II: Logical foundations

J. Algebra, **234**, 225–276, 2000.

Let A be a torsion-free abelian group and $L_{\text{gr}} = \{+, -, 0\}$ the language of abelian groups.

Axioms for $\mathbf{Qvar}(A)$ in the extended language $L_{\text{gr}A}$:

- (I) axioms of abelian groups;
- (II) torsion-free: $\forall x (nx = 0 \rightarrow x = 0)$, $n \in \mathbb{N}$;
- (III) $\forall x (p^n x = a \rightarrow x = 0)$
for all natural numbers n , all prime numbers p , and all $a \in A$,
such that the equation $p^n x = a$ has no solutions in A .

The Nullstellensatz procedure

Let $S = \{s_1 = 0, \dots, s_m = 0\}$ be a system of equations over A .

- 1** If the system S is inconsistent, then, using the Euclidean algorithm, one can use S to produce an equation of the type $p^n x = a$ with no solutions in A . In this case $\text{Rad}(S) = A[\bar{x}]$.
- 2** If the system S is consistent, then $\text{Rad}(S) = \text{Is}\langle S \rangle$.

The Nullstellensatz

for free metabelian groups



O. Chapuis

\forall -free metabelian groups

J. Symbolic Logic, **62**, 159–174, 1997.



V. Remeslennikov, R. Stöhr

On the quasivariety generated by a non-cyclic free metabelian group

Algebra Colloq., **11**, 191–214, 2004.

Elementary properties of free metabelian groups

The metabelian identity:

$$\forall x, y, z, t \quad [[x, y], [z, t]] = 1.$$

Let F_n be a free metabelian group of a rank $n \geq 2$.

- 1 The compatibility problem of systems of equations with coefficients in F_n is algorithmically undecidable (V. A. Roman'kov).
- 2 The compatibility problem of coefficient-free systems of equations over F_n is algorithmically decidable (O. Chapuis).

Fact

$\mathbf{Ucl}(F_n) = \mathbf{Ucl}(F_m)$ and $\mathbf{Qvar}(F_n) = \mathbf{Qvar}(F_m)$, $n, m \geq 2$, in the language L_{gr} . Therefore, we may talk about universal (quasi-equational) theory of a free non-abelian metabelian group and write $\mathbf{Ucl}(F)$ ($\mathbf{Qvar}(F)$).

Axioms for $\text{Ucl}(F)$ in the language L_{gr}

V. Remeslennikov, R. Stöhr:

- 1 $\forall x, y, z, t \quad [[x, y], [z, t]] = 1;$
- 2 $\forall x \quad (x^n = 1 \rightarrow x = 1), \quad n \in \mathbb{N};$
- 3 $\forall x_1 \dots, x_n, y, z \quad ([y, z]^\alpha = 1 \rightarrow [y, z] = 1),$
 $\alpha \in R_n, \varepsilon(\alpha) \neq 0;$
- 4 $\forall x_1 \dots, x_n, y, z \quad ([y, z]^\alpha = 1 \rightarrow [y, z]^{\prod(1-x_i)} = 1),$
 $\alpha \in R_n, \varepsilon(\alpha) = 0;$
- 5 CT-axiom:
 $\forall x, y, z \quad (x \neq 1 \wedge [x, y] = 1 \wedge [x, z] = 1 \rightarrow [y, z] = 1) \sim$
 $\forall x, y, z \quad ([x, y] = 1 \wedge [x, z] = 1 \rightarrow [y, z] = 1 \vee x = 1).$

Axioms for $\text{Qvar}(F)$ in the language L_{gr}

V. Remeslennikov, R. Stöhr:

$$6.1 \quad \forall x, y, z, t \quad [[x, y], [z, t]] = 1;$$

$$6.2 \quad \forall x \quad (x^n = 1 \longrightarrow x = 1), \quad n \in \mathbb{N};$$

$$6.3 \quad \forall x, y \quad ([x, y, x] = 1 \wedge [x, y, y] = 1 \longrightarrow [x, y] = 1);$$

$$6.4 \quad \forall x, y \quad ([x^n, (x^n)^y] = 1 \longrightarrow [x, x^y] = 1), \quad n \in \mathbb{N};$$

$$6.5 \quad \forall x_1 \dots, x_n, y, z \quad ([y, z]^\alpha = 1 \longrightarrow [y, z] = 1), \\ \alpha \in R_n, \varepsilon(\alpha) \neq 0;$$

$$6.6 \quad \forall x_1 \dots, x_n, y, z \quad ([y, z]^\alpha = 1 \longrightarrow [y, z]^{\prod(1-x_i)} = 1), \\ \alpha \in R_n, \varepsilon(\alpha) = 0;$$

$$6.7 \quad \forall x_1 \dots, x_n, y, z \quad ([y, z]^\alpha = 1 \longrightarrow [y, z]^{\beta_i} = 1), \\ \text{Rad}(\alpha) = \langle \beta_1, \dots, \beta_q \rangle;$$

$$6.8 \quad \forall x_1 \dots, x_n, y_1, \dots, y_n, z \\ (z = \prod [x_i, y_i] \wedge \bigwedge_{i=1}^n [z, x_i] = [z, y_i] = 1 \longrightarrow z = 1).$$

Proposition 3.2

There exists an algorithm to compute the radical $\text{Rad}(S)$ of a finite coefficient-free system of equations S over a free finitely generated non-abelian metabelian group.

O. Chapuis's Theorem

For any finitely generated group G the following conditions are equivalent:

- G is a subgroup of a wreath product $W_{r,s}$, $r, s \in \mathbb{N}$;
- $G \in \mathbf{Ucl}(F)$;
- G satisfies 6.1, 6.2, 6.5, 6.6, and CT;
- G is an s -group.

Questions and problems

Open questions

- 1 Does there exist a satisfactory system of axioms for $\mathbf{Ucl}(\mathbb{Q}_p)$ and $\mathbf{Qvar}(\mathbb{Q}_p)$ in the language $L_{\text{ring}\mathbb{Q}_p}$? And for $\mathbf{Qvar}(\mathbb{Q}_p)$ in the language $L_{\text{ring}} \cup \{x \mid_p y\}$?
- 2 Let F be a free solvable group of class $r > 2$. Does there exist a recursive system of axioms for $\mathbf{Ucl}(F)$ and $\mathbf{Qvar}(F)$ in the language L_{gr} ? In other languages?
- 3 Nullstellensatz for free metabelian (free solvable) group in the extended language
 - in dimension 1,
 - for semigroup equations.