

Variations on Pebbling and Graham's Conjecture

David S. Herscovici
Quinnipiac University
with Glenn H. Hurlbert
and Ben D. Hester
Arizona State University

Talk structure

Pebbling numbers

Products and Graham's Conjecture

Variations

- Optimal pebbling

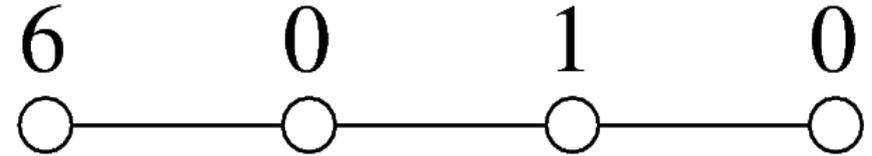
- Weighted graphs

- Choosing target distributions

Basic notions

Distributions on G :

$$D: V(G) \rightarrow \mathbb{N}$$



$D(v)$ counts pebbles
on v

Pebbling moves

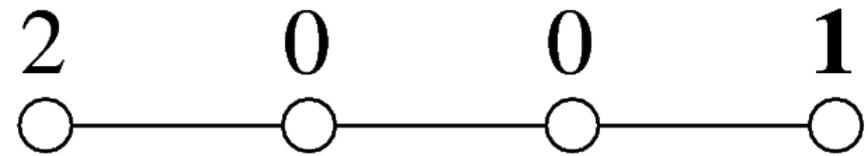
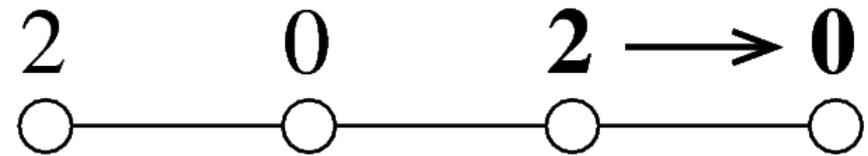
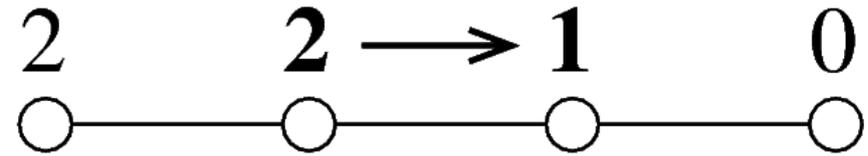
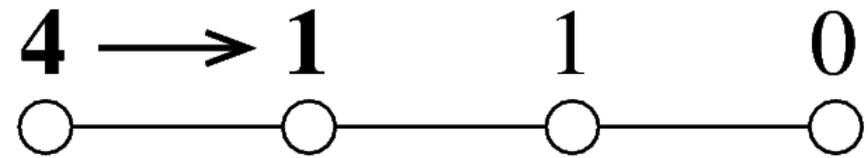
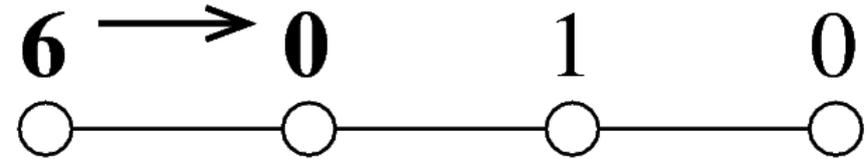
Basic notions

Distributions on G :

$$D: V(G) \rightarrow \mathbb{N}$$

$D(v)$ counts pebbles
on v

Pebbling moves



Pebbling numbers

$\pi(G, D)$ is the number of pebbles required to ensure that D can be reached from any distribution of $\pi(G, D)$ pebbles.

If S is a **set** of distributions on G

$$\pi(G, S) = \max_{D \in S} \pi(G, D)$$

Pebbling numbers

$\pi(G, D)$ is the number of pebbles required to ensure that D can be reached from any distribution of $\pi(G, D)$ pebbles.

If S is a **set** of distributions on G

$$\pi(G, S) = \max_{D \in S} \pi(G, D)$$

Optimal pebbling number:

$\pi^*(G, S)$ is the number of pebbles required in **some** distribution from which every $D \in S$ can be reached

Common pebbling numbers

$S_1(G)$: 1 pebble anywhere

$$\pi(G) = \pi(G, S_1(G))$$

(pebbling number)

$S_t(G)$: t pebbles on some vertex

$$\pi_t(G) = \pi(G, S_t(G))$$

(t -pebbling number)

δ_v : One pebble on v

$$\pi(G, v) = \pi(G, \delta_v)$$

$S(G, t)$ and $\pi(G, t)$

$S(G, t)$: all distributions with a total of t pebbles (anywhere on the graph)

Conjecture

$$\pi(G, S(G, t)) = \pi(G, S_t(G))$$

i.e. hardest-to-reach ***target*** configurations have all pebbles on one vertex

True for K_n , C_n , trees

Cover pebbling

$\Gamma(G)$: one pebble on every vertex

$$\gamma(G) = \pi(G, \Gamma(G))$$

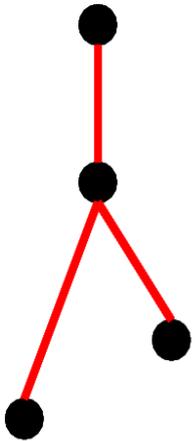
(*cover pebbling number*)

Sjöstrund: If $D(v) \geq 1$ for all vertices v , then there is a critical distribution with all pebbles on one vertex.

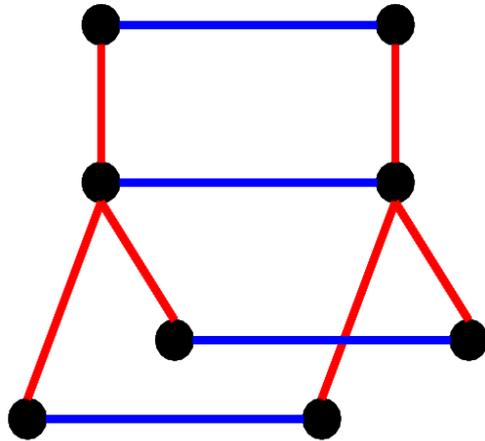
(A *critical* distribution has one pebble less than the required number and cannot reach some target distribution)

Cartesian products of graphs

K_2



G



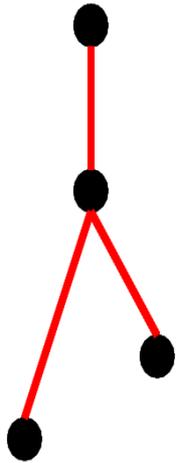
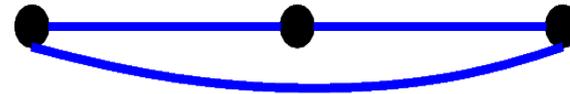
$G \times K_2$

Cartesian products of graphs

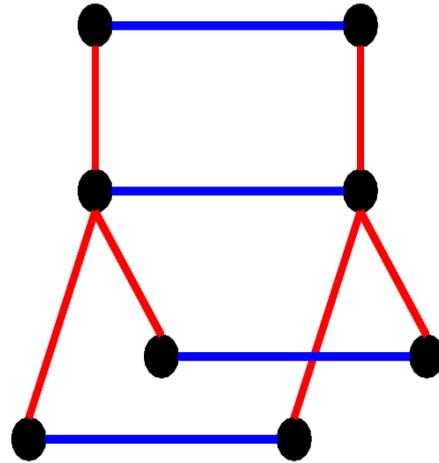
K_2



K_3



G



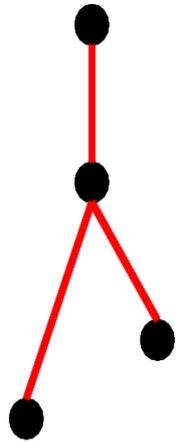
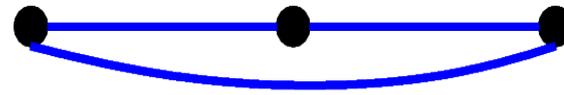
$G \times K_2$

Cartesian products of graphs

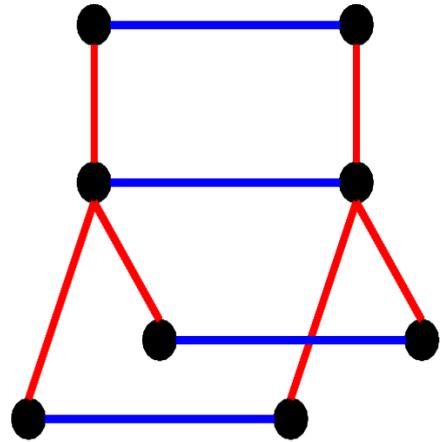
K_2



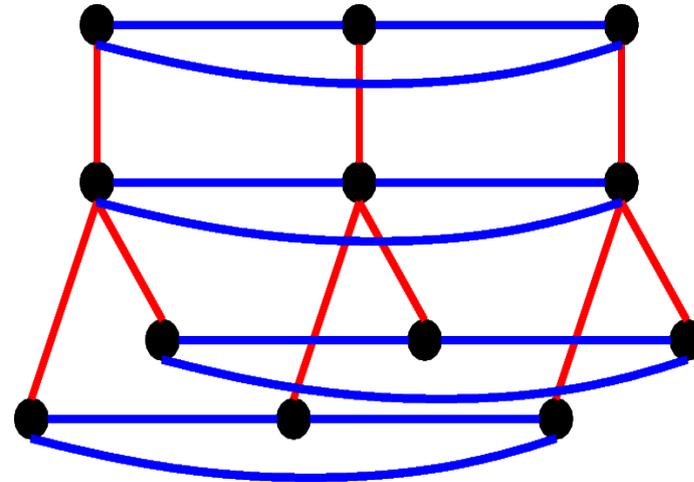
K_3



G



$G \times K_2$



$G \times K_3$

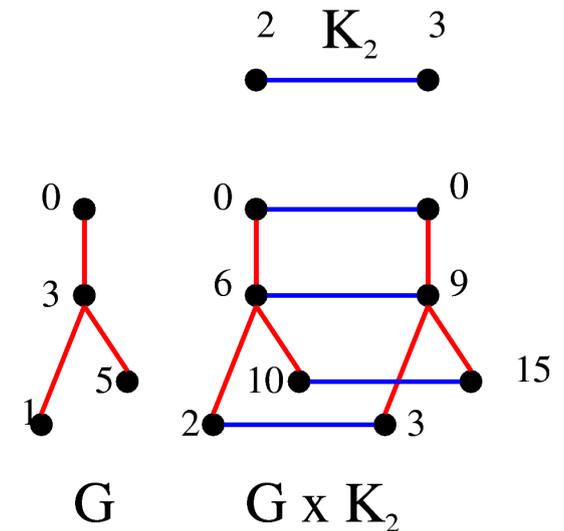
Products of distributions

Product of distributions:

D_1 on G ; D_2 on H

then $D_1 \cdot D_2$ on $G \times H$

$$D_1 \cdot D_2((v, w)) = D_1(v) D_2(w)$$



Products of distributions

Product of distributions:

D_1 on G ; D_2 on H

then $D_1 \cdot D_2$ on $G \times H$

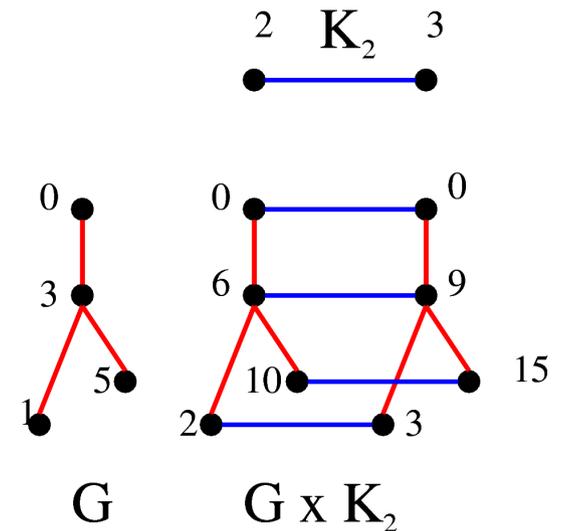
$$D_1 \cdot D_2((v, w)) = D_1(v) D_2(w)$$

Products of **sets** of distributions

S_1 a set of distros on G

S_2 a set of distros on H

$$S_1 \cdot S_2 = \{D_1 \cdot D_2 : D_1 \in S_1 ; D_2 \in S_2\}$$



Graham's Conjecture generalized

Graham's Conjecture:

$$\pi(G \times H) \leq \pi(G)\pi(H)$$

Generalization:

$$\pi(G \times H, S_1 \cdot S_2) \leq \pi(G, S_1)\pi(H, S_2)$$

Optimal pebbling

Observation (not obvious):

If we can get from D_1 to D_1' in G and from D_2 to D_2' in H , we can get from $D_1 \cdot D_2$ to $D_1 \cdot D_2'$ in $G \times H$ to $D_1' \cdot D_2'$ in $G \times H$

Conclusion (optimal pebbling):

$$\pi^*(G \times H, S_1 \cdot S_2) \leq \pi^*(G, S_1) \pi^*(H, S_2)$$

Graham's Conjecture holds for optimal pebbling in most general setting

Weighted graphs

Edges have weights w

Pebbling moves: remove w pebbles from one vertex, move 1 to adjacent vertex

$\pi(G)$, $\pi(G, D)$ and $\pi(G, S)$ still make sense

$G \times H$ also makes sense:

$$\text{wt}(\{(v, w), (v, w')\}) = \text{wt}(\{w, w'\})$$

$$\text{wt}(\{(v, w), (v', w)\}) = \text{wt}(\{v, v'\})$$

The Good news

Chung: Hypercubes ($K_2 \times K_2 \times \dots \times K_2$) satisfies Graham's Conjecture for any collection of weights on the edges

We can focus on complete graphs in most applications

The Bad News

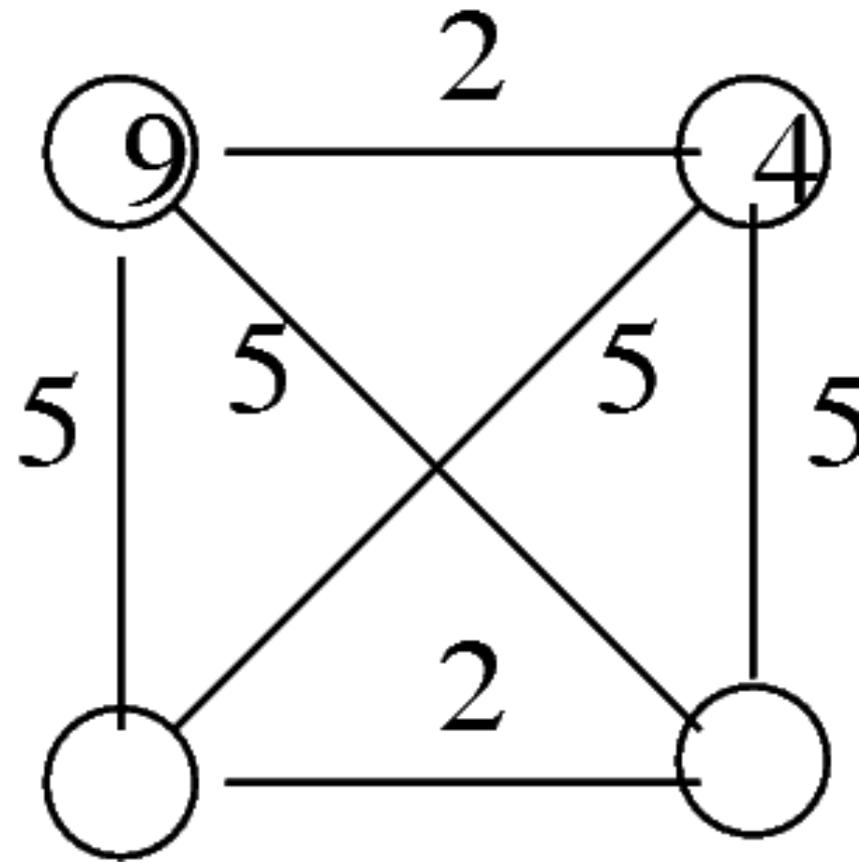
Complete graphs are
hard!

The Bad News

Complete graphs are hard!

Sjöstrand's Theorem fails

13 pebbles on one vertex can cover K_4 , but...



Some specializations

Conjecture 1:

$$\pi(G \times H, (v, w)) \leq \pi(G, v) \pi(H, w)$$

Conjecture 2:

$$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$$

Clearly Conjecture 2 implies Conjecture 1

Some specializations

Conjecture 1:

$$\pi(G \times H, (v, w)) \leq \pi(G, v) \pi(H, w)$$

Conjecture 2:

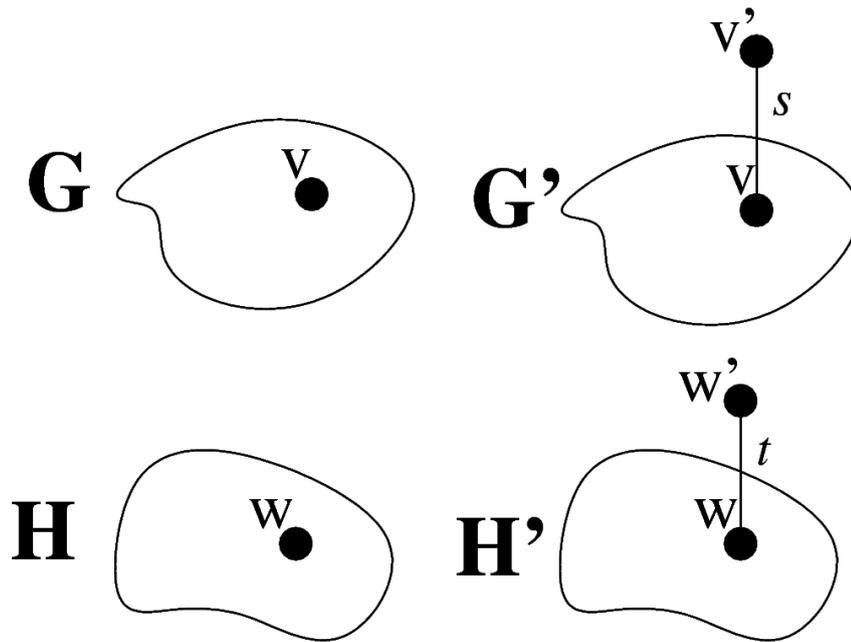
$$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$$

Clearly Conjecture 2 implies Conjecture 1

These conjectures are equivalent on weighted graphs

$\pi(G \times H, (v, w)) \leq \pi(G, v) \pi(H, w)$
 implies

$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$



$$\pi(G', v') = \pi_s(G, v)$$

$$\pi(H', w') = \pi_t(H, w)$$

$$\pi(G \times H, (v, w)) \leq \pi(G, v) \pi(H, w)$$

implies

$$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$$

If st pebbles cannot be moved to (v, w) from D in $G \times H$, then (v', w') cannot be reached from D in $G' \times H'$ (delay moves onto $\{v'\} \times H'$ and $G' \times \{w'\}$ as long as possible)

$$\begin{aligned} \pi_{st}(G \times H, (v, w)) &\leq \pi(G' \times H', (v', w')) \\ &\leq \pi(G', v') \pi(H', w') = \pi_s(G, v) \pi_t(H, w) \end{aligned}$$

Implications for regular pebbling

Conjecture 1:

$$\pi(G \times H, (v, w)) \leq \pi(G, v) \pi(H, w)$$

equivalent to Conjecture 1':

$$\pi_{2^{ab}}(G \times H, (v, w)) \leq \pi_{2^a}(G, v) \pi_{2^b}(H, w)$$

Implications for regular pebbling

Conjecture 1:

$$\pi(G \times H, (v, w)) \leq \pi(G, v) \pi(H, w)$$

equivalent to Conjecture 1':

$$\pi_{2^{ab}}(G \times H, (v, w)) \leq \pi_{2^a}(G, v) \pi_{2^b}(H, w)$$

Conjecture 2:

$$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$$

equivalent to Conjecture 2':

$$\pi_{st}(G \times H, (v, w)) \leq \pi_s(G, v) \pi_t(H, w)$$

when s and t are odd

Choosing a target

Observation: To reach an unoccupied vertex v in G , we need to put two pebbles on any neighbor of v .

We can choose the target neighbor

If S is a *set* of distributions on G , $\rho(G, S)$ is the number of pebbles needed to reach ***some*** distribution in S

Idea: Develop an induction argument to prove Graham's conjecture

Comparing pebbling numbers

$\forall D \in \mathcal{S} \forall D' \in \mathcal{S}(G, \pi(G, S))$

D is reachable from D' by a sequence of pebbling moves

Comparing pebbling numbers

$$\forall D \in S \forall D' \in S(G, \pi(G, S))$$

D is reachable from D' by a sequence of pebbling moves

$$\forall D \in S \exists D' \in S(G, \pi^*(G, S))$$

D is reachable from D' by a sequence of pebbling moves

Comparing pebbling numbers

$$\forall D \in S \forall D' \in S(G, \pi(G, S))$$

D is reachable from D' by a sequence of pebbling moves

$$\forall D \in S \exists D' \in S(G, \pi^*(G, S))$$

D is reachable from D' by a sequence of pebbling moves

$$\forall D' \in S(G, \rho(G, S)) \exists D \in S$$

D is reachable from D' by a sequence of pebbling moves

Properties of $\rho(G, S)$

$$\rho_1(G) = 1$$

$$\rho_2(G) = |V(G)| + 1$$

Properties of $\rho(G, S)$

$$\rho_1(G) = 1$$

$$\rho_2(G) = |V(G)| + 1$$

SURPRISE!

Graham's Conjecture fails!

Let H be the trivial graph: $S_H = \{2\delta_v\}$

$$\rho(G \times H, S_1(G) \cdot S_H) = \rho_2(G) = |V(G)| + 1$$

$$\rho(G, S_1(G)) = \rho_1(G) = 1$$

$$\rho(H, S_H) = 2$$